

# **Bounds on quantum process fidelity from quantum state fidelity measurements**

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# Quantum operations

Choi-Jamiolkowski isomorphism

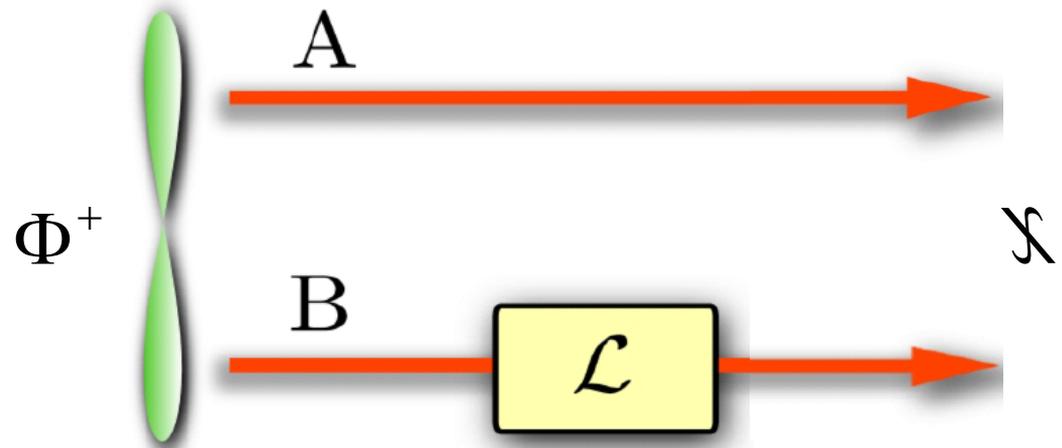
$$\chi = I \otimes L(\Phi^+)$$

Maximally entangled probe state

$$|\Phi^+\rangle = \sum_{j=1}^d |jj\rangle$$

Transformation of input states

$$\rho_{\text{out}} = \text{Tr}_{\text{in}}[\rho_{\text{in}}^T \otimes \mathbb{I}_{\text{out}} \chi]$$



Complete positivity condition

$$\chi \geq 0$$

Trace preservation condition

$$\text{Tr}_{\text{out}}[\chi] = \mathbb{I}_{\text{in}}$$

# Quantum gate fidelity

Choi matrix of a unitary operation  $U$  – density matrix of a pure maximally entangled state:

$$\chi_U = I \otimes U |\Phi^+\rangle\langle\Phi^+| I \otimes U^\dagger$$

Quantum gate fidelity – defined as normalized overlap of Choi operators:

$$F_\chi = \frac{\text{Tr}[\chi_U \chi]}{\text{Tr}[\chi_U] \text{Tr}[\chi]}$$

Practical determination of gate fidelity:

- quantum process tomography – full reconstruction of process matrix  $\chi$
- Monte Carlo sampling
- Hofmann lower and upper bounds on gate fidelity

# Hofmann bounds on gate fidelity

Determine average state fidelities for two conjugate bases  $|\psi_j\rangle$   $|\varphi_k\rangle$

$$\langle \psi_k | \psi_j \rangle = \delta_{jk} \quad \langle \varphi_k | \varphi_j \rangle = \delta_{jk} \quad |\langle \varphi_k | \psi_j \rangle| = \frac{1}{d}$$

**Average state fidelities:**

$$F_1 = \frac{1}{d} \sum_{j=1}^d \langle \psi_j | \rho_{j,out} | \psi_j \rangle$$

$$F_2 = \frac{1}{d} \sum_{k=1}^d \langle \varphi_k | \rho_{k,out} | \varphi_k \rangle$$

$$\rho_{j,out} = \text{Tr}_{in} \left[ |\psi_j\rangle \langle \psi_j| \otimes I \chi \right]$$

$$\rho_{k,out} = \text{Tr}_{in} \left[ |\varphi_k\rangle \langle \varphi_k| \otimes I \chi \right]$$

# Hofmann bounds on gate fidelity

Average state fidelities:

$$F_1 = \frac{1}{d} \sum_{j=1}^d \langle \psi_j | \rho_{j,\text{out}} | \psi_j \rangle$$

$$F_2 = \frac{1}{d} \sum_{k=1}^d \langle \varphi_k | \rho_{k,\text{out}} | \varphi_k \rangle$$

$$\rho_{j,\text{out}} = \text{Tr}_{\text{in}} \left[ |\psi_j\rangle\langle\psi_j| \otimes I_{\chi} \right]$$

$$\rho_{k,\text{out}} = \text{Tr}_{\text{in}} \left[ |\varphi_k\rangle\langle\varphi_k| \otimes I_{\chi} \right]$$

Lower bound on quantum gate fidelity

$$F_{\chi} \geq F_1 + F_2 - 1$$

Upper bound on quantum gate fidelity

$$F_{\chi} \leq \min(F_1, F_2)$$

# Minimum number of probe states

To obtain a nontrivial bound on quantum gate fidelity, it suffices to probe the quantum gate with  $d+1$  pure probe states:

**Computational basis states:**

$$|j\rangle, j = 0, \dots, d-1$$

**Superposition state:**

$$|+\rangle = \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle$$

**Average state fidelity F:**

$$F = \frac{1}{d} \text{Tr} [|j\rangle\langle j| \otimes |j\rangle\langle j| \chi]$$

**State fidelity G:**

$$G = \text{Tr} [|+\rangle\langle +| \otimes |+\rangle\langle +| \chi]$$

**Determine a lower bound on quantum gate fidelity given F and G.**

# Two-qubit gates

Construction of a specific quantum operation that will be proven to minimize the gate fidelity for given fixed state fidelities  $F$  and  $G$ :

$$\tilde{\chi} = (\mathbb{I} \otimes U) \tilde{\chi}_S (\mathbb{I} \otimes U^\dagger)$$

$$\tilde{\chi}_S = \sum_{m=0}^3 |\chi_m\rangle \langle \chi_m|$$

$$|\chi_0\rangle = a Z_{00} |\Phi_2^+\rangle + b |++\rangle |++\rangle$$

$$|\chi_1\rangle = c Z_{01} |\Phi_2^+\rangle + d |++\rangle |+-\rangle$$

$$|\chi_2\rangle = c Z_{10} |\Phi_2^+\rangle + d |++\rangle |-+\rangle$$

$$|\chi_3\rangle = c Z_{11} |\Phi_2^+\rangle + d |++\rangle |--\rangle$$

# Parameters of the quantum operation

$$\tilde{\chi}_S = \sum_{m=0}^3 |\chi_m\rangle\langle\chi_m|$$

$$|\chi_0\rangle = a Z_{00}|\Phi_2^+\rangle + b |++\rangle|++\rangle$$

$$|\chi_1\rangle = c Z_{01}|\Phi_2^+\rangle + d |++\rangle|+-\rangle$$

$$|\chi_2\rangle = c Z_{10}|\Phi_2^+\rangle + d |++\rangle| -+\rangle$$

$$|\chi_3\rangle = c Z_{11}|\Phi_2^+\rangle + d |++\rangle|--\rangle$$

**Determined from the trace-preservation condition and from the fixed state fidelities:**

$$a = \frac{2}{3} \left[ (8F - 5)\sqrt{G} - 4\sqrt{(1 - F)(4F - 1)(1 - G)} \right]$$

$$b = \sqrt{G} - \frac{a}{2},$$

$$c = \sqrt{\frac{4 - a^2}{3}},$$

$$d = \sqrt{\frac{1 - G}{3}} - \frac{1}{2}\sqrt{\frac{4 - a^2}{3}}.$$

**Analytical formula for quantum gate fidelity of this operation:**

$$\tilde{F}_\chi = \left[ (2F - 1)\sqrt{G} - \sqrt{(4F - 1)(1 - F)}\sqrt{1 - G} \right]^2$$

# Generalized Hofmann lower bound

Define a threshold fidelity

$$F_{\text{th}} = \frac{1}{8} \left( 5 - G + \sqrt{9 - 10G + G^2} \right)$$

If  $F > F_{\text{th}}$  then the lower bound on gate fidelity reads

$$\tilde{F}_\chi = \left[ (2F - 1)\sqrt{G} - \sqrt{(4F - 1)(1 - F)}\sqrt{1 - G} \right]^2$$

If  $F < F_{\text{th}}$  then the lower bound is zero:

$$F_\chi \geq 0$$

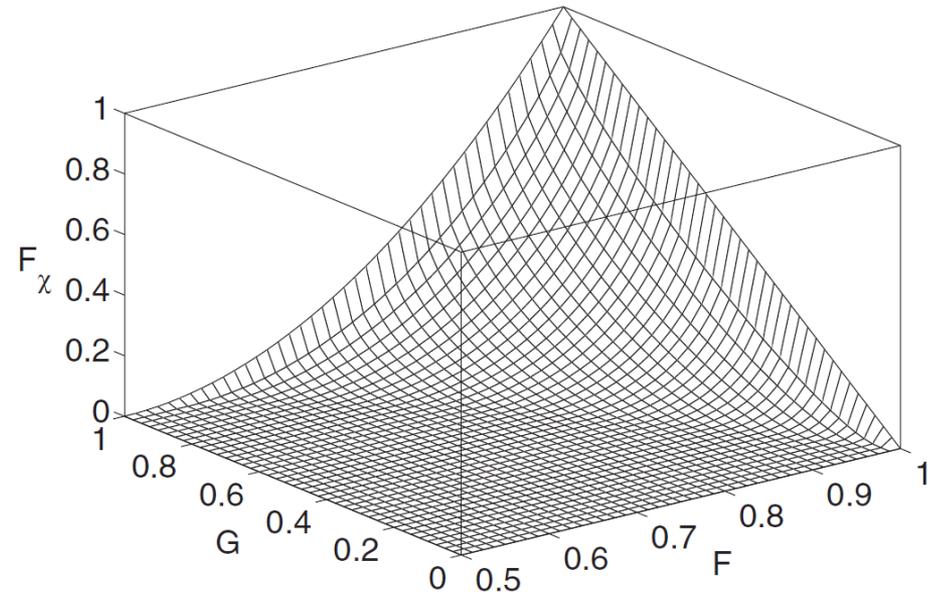


FIG. 1. Lower bound on quantum process fidelity  $F_\chi$  of a two-qubit quantum operation plotted as a function of state fidelities  $F$  and  $G$ .

# Comparison with standard Hofmann bound

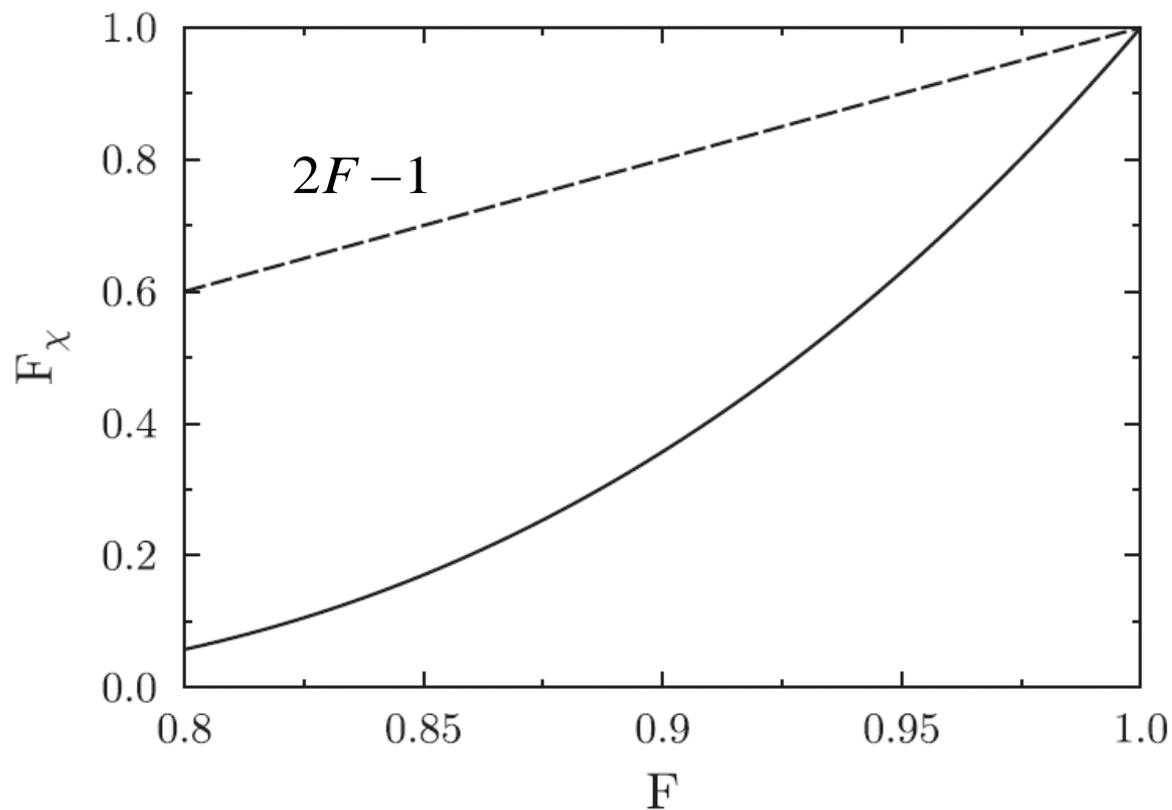


FIG. 2. Lower bound on quantum process fidelity  $F_\chi$  determined from the knowledge of state fidelities  $F$  and  $G$  (solid line) and the Hofmann lower bound on quantum process fidelity (dashed line) plotted for two-qubit quantum operations assuming  $G = F' = F$ .

# Optimality proof

The proof exploits semidefinite programming techniques.

Define operator

$$M = \frac{1}{4} |\Phi_2^+\rangle \langle \Phi_2^+| + x R_F + w R_G + y \mathbb{I} + z |++\rangle \langle ++| \otimes \mathbb{I}_{\text{out}}$$

where

$$R_F = \frac{1}{4} \sum_{j,k=0}^1 |jk\rangle \langle jk| \otimes |jk\rangle \langle jk| \quad R_G = |++\rangle \langle ++| \otimes |++\rangle \langle ++|$$

and find Lagrange multipliers  $x, w, y, z$  such that

$$M \geq 0$$

$$M \tilde{\chi}_S = 0$$

# Optimality proof II

$$M = \frac{1}{4}|\Phi_2^+\rangle\langle\Phi_2^+| + xR_F + wR_G + y\mathbb{I} + z|++\rangle\langle++| \otimes \mathbb{I}_{\text{out}}$$

Positive semidefiniteness of  $M$  implies

$$\text{Tr}[M\chi] \geq 0 \quad \longrightarrow \quad F_\chi + xF + wG + 4y + z \geq 0$$

$$M\tilde{\chi}_S = 0 \quad \longrightarrow \quad \tilde{F}_\chi = -xF - wG - 4y - z$$

Altogether we have  $F_\chi \geq \tilde{F}_\chi$

# Optimality proof III

$$M = \frac{1}{4} |\Phi_2^+\rangle \langle \Phi_2^+| + x R_F + w R_G + y \mathbb{I} + z |+++ \rangle \langle +++| \otimes \mathbb{I}_{\text{out}}$$

**Lagrange multipliers:**

$$x = \frac{\sqrt{4-a^2} (3a + 2\sqrt{G})}{2a\sqrt{1-G} - 2\sqrt{G(4-a^2)}},$$

$$w = - \left( 3\sqrt{4-a^2} + 2\sqrt{1-G} \right) \frac{3a + 2\sqrt{G}}{64\sqrt{G(1-G)}},$$

$$y = \frac{1}{32} \left( 3a + 2\sqrt{G} \right) \frac{3\sqrt{4-a^2} + 2\sqrt{1-G}}{\sqrt{G(4-a^2)} - a\sqrt{1-G}},$$

$$z = \frac{\sqrt{4-a^2} - 2\sqrt{1-G}}{2\sqrt{1-G}} y.$$

$$a = \frac{2}{3} \left[ (8F - 5)\sqrt{G} - 4\sqrt{(1-F)(4F-1)(1-G)} \right]$$

**Eigenvalues of M:**

$$\lambda_1 = y,$$

$$\lambda_2 = \frac{1}{8}(A - \sqrt{B}), \quad \lambda_3 = \frac{1}{8}(A + \sqrt{B}),$$

$$\lambda_4 = \frac{1}{8}(C - \sqrt{D}), \quad \lambda_5 = \frac{1}{8}(C + \sqrt{D}).$$

$$A = x + 8y + 4z, \quad B = x^2 - 4xz + 16z^2,$$

$$C = 1 + 4w + x + 8y + 4z,$$

$$D = 1 + 16w^2 + 2x + x^2 - 4w(1 + x - 8z) - 4z - 4xz + 16z^2.$$

It can be proven analytically that all the eigenvalues are non-negative.

# N-qubit gates

Construction of the specific quantum operation can be extended to N-qubit gates but no proof of optimality is available.

$$\tilde{\chi}_s = \sum_{j=0}^{2^N-1} |\chi_j\rangle\langle\chi_j|$$

$$|\chi_0\rangle = a|\Phi_N^+\rangle + b|s\rangle|s\rangle$$

$$|\chi_j\rangle = \mathbb{I} \otimes V_j(c|\Phi_N^+\rangle + d|s\rangle|s\rangle)$$

$$|\Phi_N^+\rangle = \frac{1}{\sqrt{2^N}} \sum_{j=0}^{2^N-1} |j\rangle|j\rangle \quad V_j = \bigotimes_{k=1}^N \sigma_Z^{j_k}$$

Anyway, this construction yields an upper bound on the generalized Hofmann lower bound for this case:

$$\tilde{F}_\chi = \left\{ [1 - (1 - F)2^{N-1}] \sqrt{G} - \sqrt{(1 - F)(1 - G)} \sqrt{2^N - 1 - (1 - F)2^{2N-2}} \right\}^2$$

$$F_{\text{th}} = 1 - \frac{1}{2^{N-1}} + \frac{1 - G}{2^{2N-1}} + \frac{2}{2^{2N}} \sqrt{(1 - G)[(2^N - 1)^2 - G]}$$

The bound is 0 when  $F < F_{\text{th}}$ .

The fidelity F must be exponentially close to 1 to obtain a nontrivial bound.