

Compensation of side-channel noise infusion on the receiver side in continuous-variable quantum key distribution

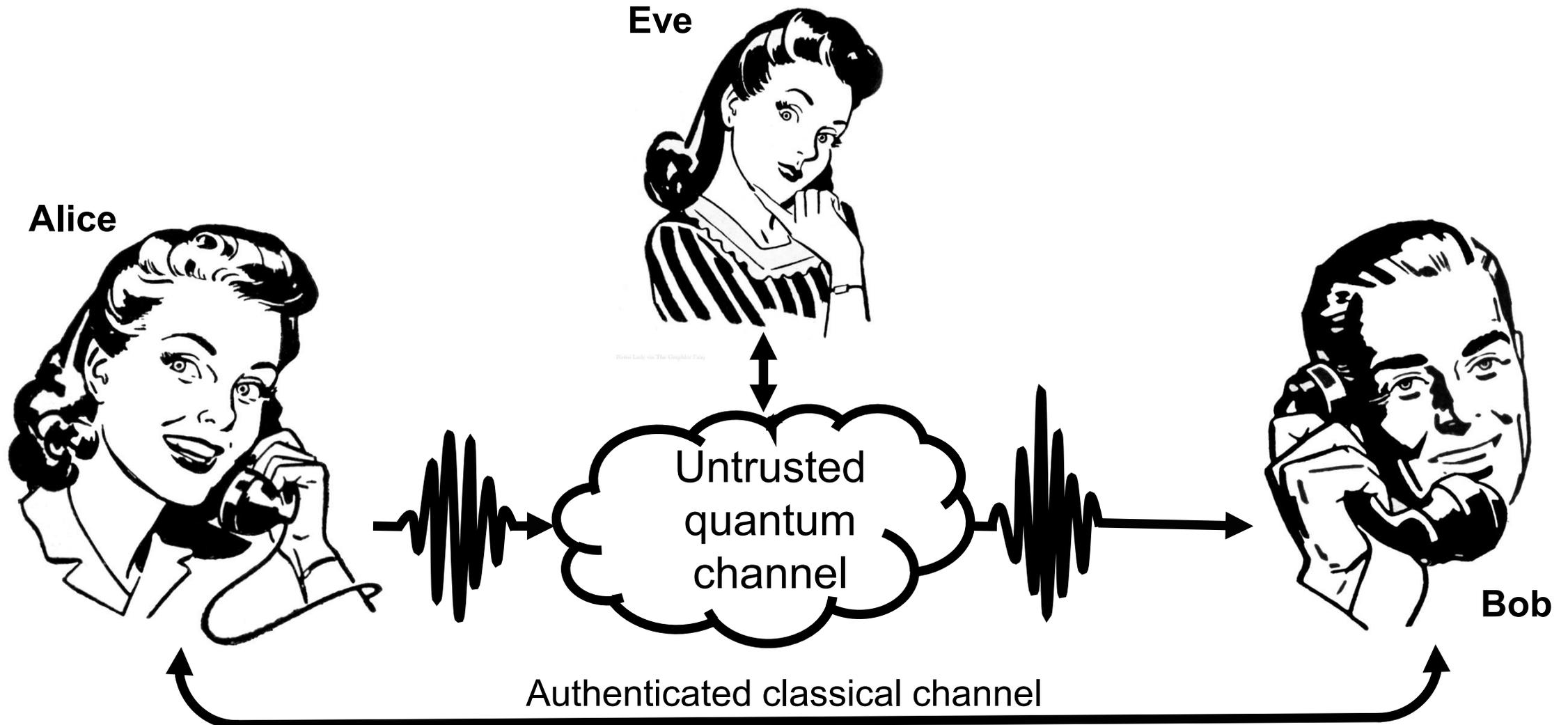
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Outline

- Motivation
- Definition and types of side channels
- Side channel negative effects
- Decoupling of information leakage on sender side
- Decoupling of noise infusion on receiver side
- Summary

Quantum key distribution



Imperfections

- Trusted preparation noise can break the security in reverse reconciliation protocol [1], but can be purified [2] or tolerated in the direct reconciliation scheme [3, 4]

[15] R. Filip, Phys. Rev. A 77, 022310 (2008).

[16] V. C. Usenko, and R. Filip Phys. Rev. A 81, 022318 (2010).

[17] C. Weedbrook, S. Pirandola, S. Lloyd, and T. C. Ralph, Phys. Rev. Lett. 105, 110501 (2010)

[18] C. Weedbrook, S. Pirandola, T. C. Ralph Phys. Rev. A 86, 022318 (2012)

Imperfections

- Trusted preparation noise can break the security in reverse reconciliation protocol [1], but can be purified [2] or tolerated in the direct reconciliation scheme [3, 4]
- Trusted detection noise limits the key rate, but can be partially helpful to make the protocol more robust against the noise in the quantum channel [5]

[15] R. Filip, Phys. Rev. A 77, 022310 (2008).

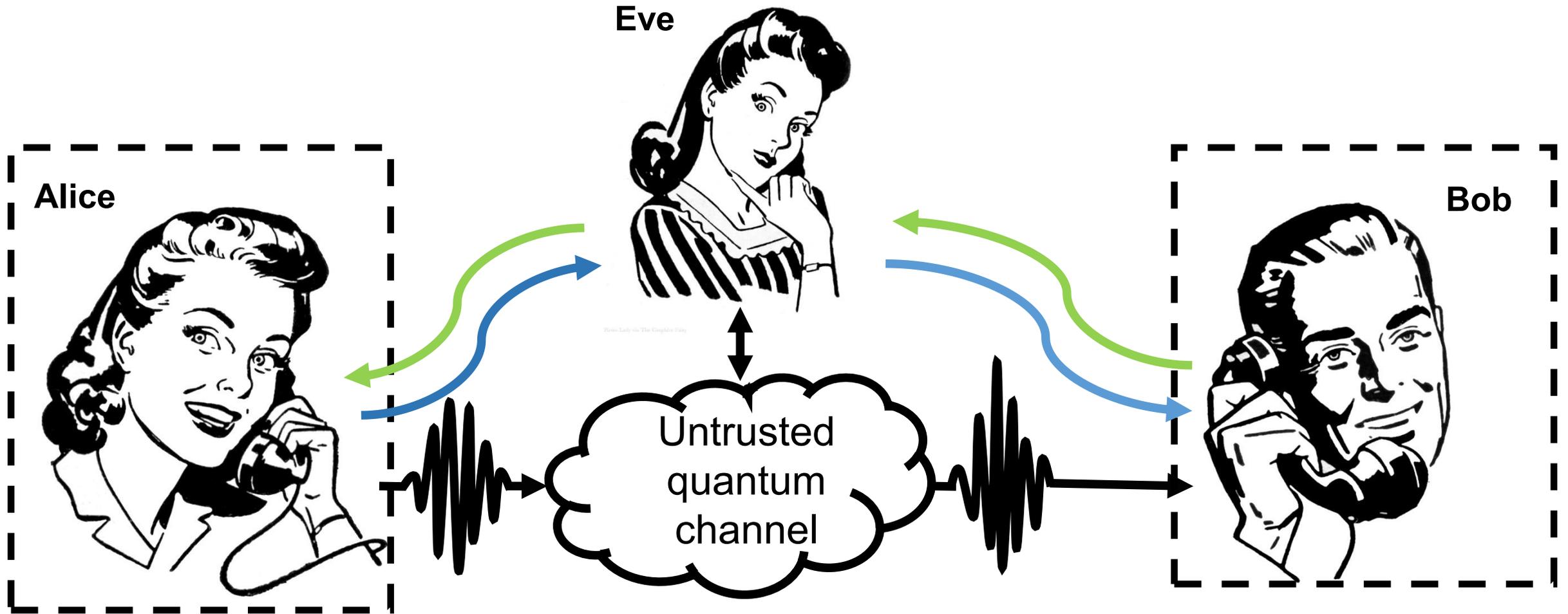
[16] V. C. Usenko, and R. Filip Phys. Rev. A 81, 022318 (2010).

[17] C. Weedbrook, S. Pirandola, S. Lloyd, and T. C. Ralph, Phys. Rev. Lett. 105, 110501 (2010)

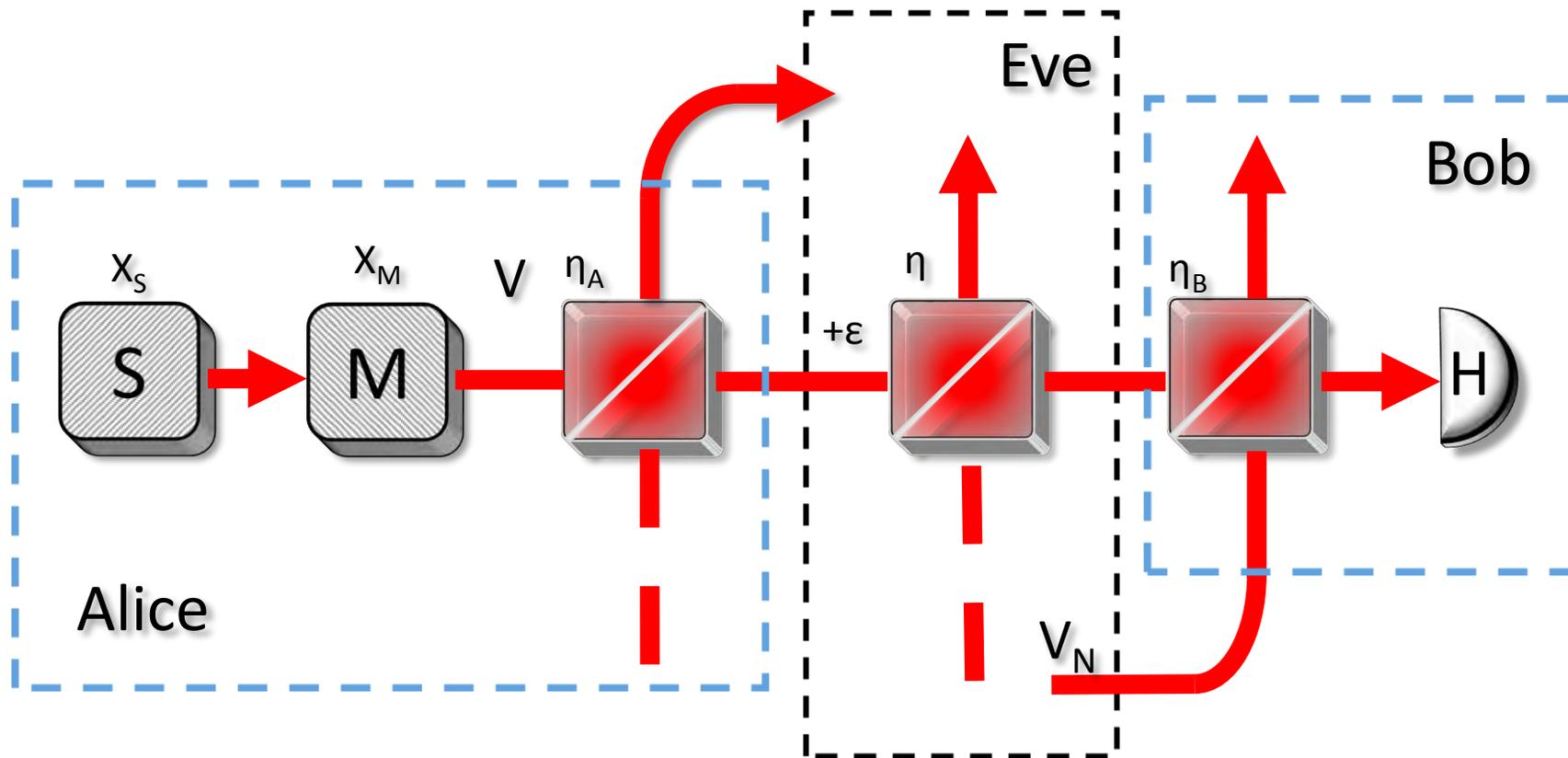
[18] C. Weedbrook, S. Pirandola, T. C. Ralph Phys. Rev. A 86, 022318 (2012)

[5] R. García-Patón, N. J. Cerf, Phys. Rev. Lett. 102, 130501 (2009)

Quantum key distribution



Side channels



S – source
(laser/OPO)
M – quadrature
modulator
 η - untrusted
channel loss
 ϵ – untrusted
channel excess
noise
H – homodyne
detector

Side channels negative effect

Key rate for individual attacks:

$$K_{ind} = \beta I_{AB} - I_{BE}$$

Sender-side loss

$$I_{AB} = \frac{1}{2} \log_2 \frac{1}{1 - \frac{\eta_A \eta V_M}{\eta_A \eta (V-1) + 1}}$$

$$I_{BE} = \frac{1}{2} \log_2 \frac{[\eta_A \eta (V-1) + 1][V - \eta_A \eta (V-1)]}{V}$$

$$K_{V \rightarrow \infty} = \log_2 \frac{1}{1 - \eta_A \eta}$$

Key rate for collective attacks:

$$K_{col} = \beta I_{AB} - \chi_{BE}$$

Receiver-side noise infusion

$$I_{AB} = \frac{1}{2} \log_2 \frac{1}{1 - \frac{\eta_B \eta V_M}{\eta_B (\eta V + 1 - \eta) + (1 - \eta_B) V_N}}$$

$$I_{BE} = \frac{1}{2} \log_2 \frac{\eta_B (\eta V + 1 - \eta) + (1 - \eta_B) V_N}{\frac{\eta_B V}{\eta + (1 - \eta) V} + \frac{1 - \eta_B}{V_N}}$$

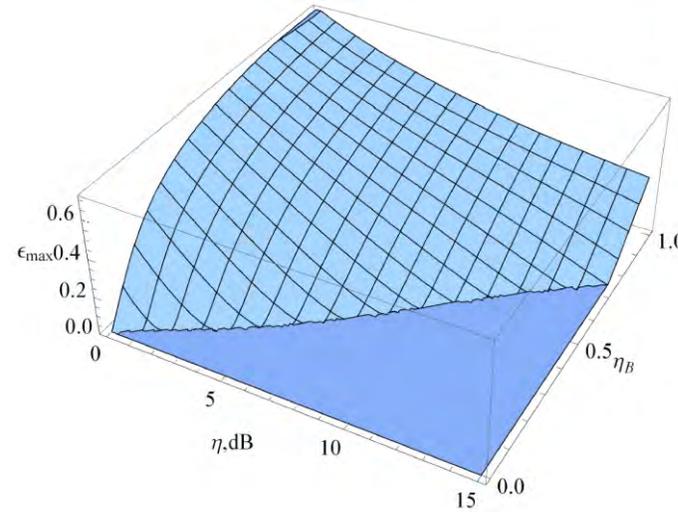
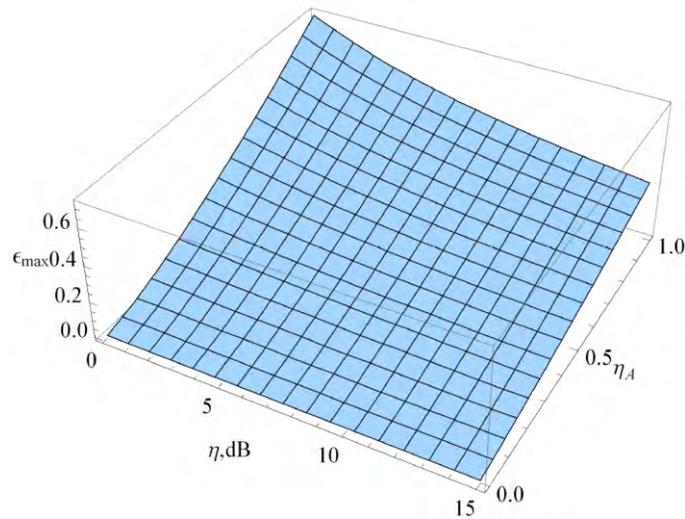
$$V_N^{max} \Big|_{\eta \rightarrow 0}^{V \rightarrow \infty} = \frac{1}{1 - \eta_B}$$

Side channels negative effect

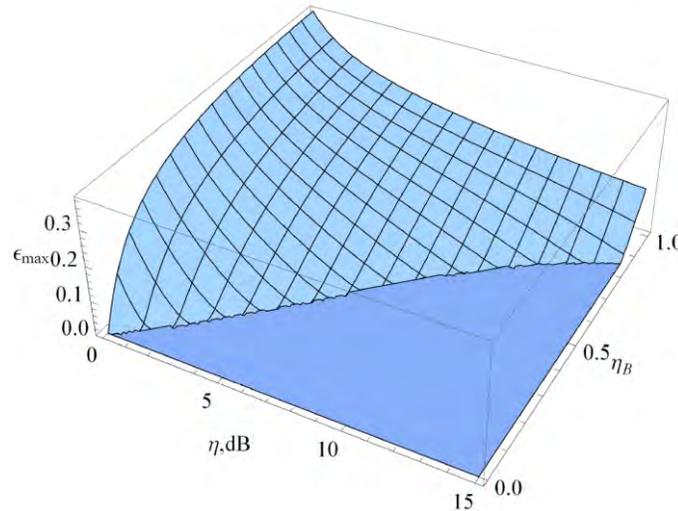
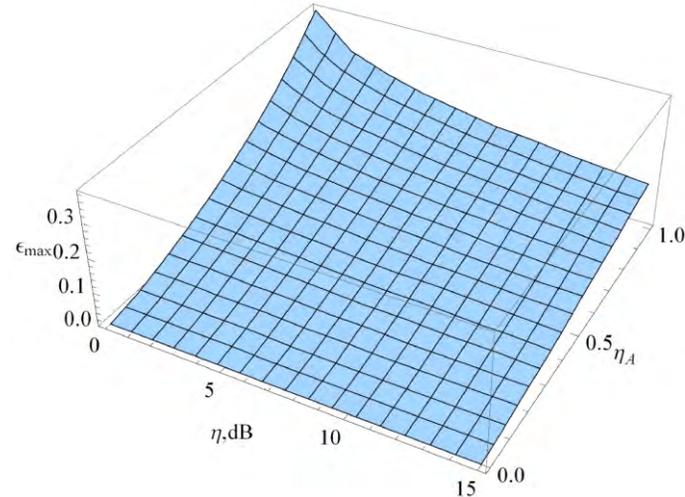
Leakage on the sender side

Noise infusion on the receiver side

Squeezed-state
protocol



Coherent-state
protocol

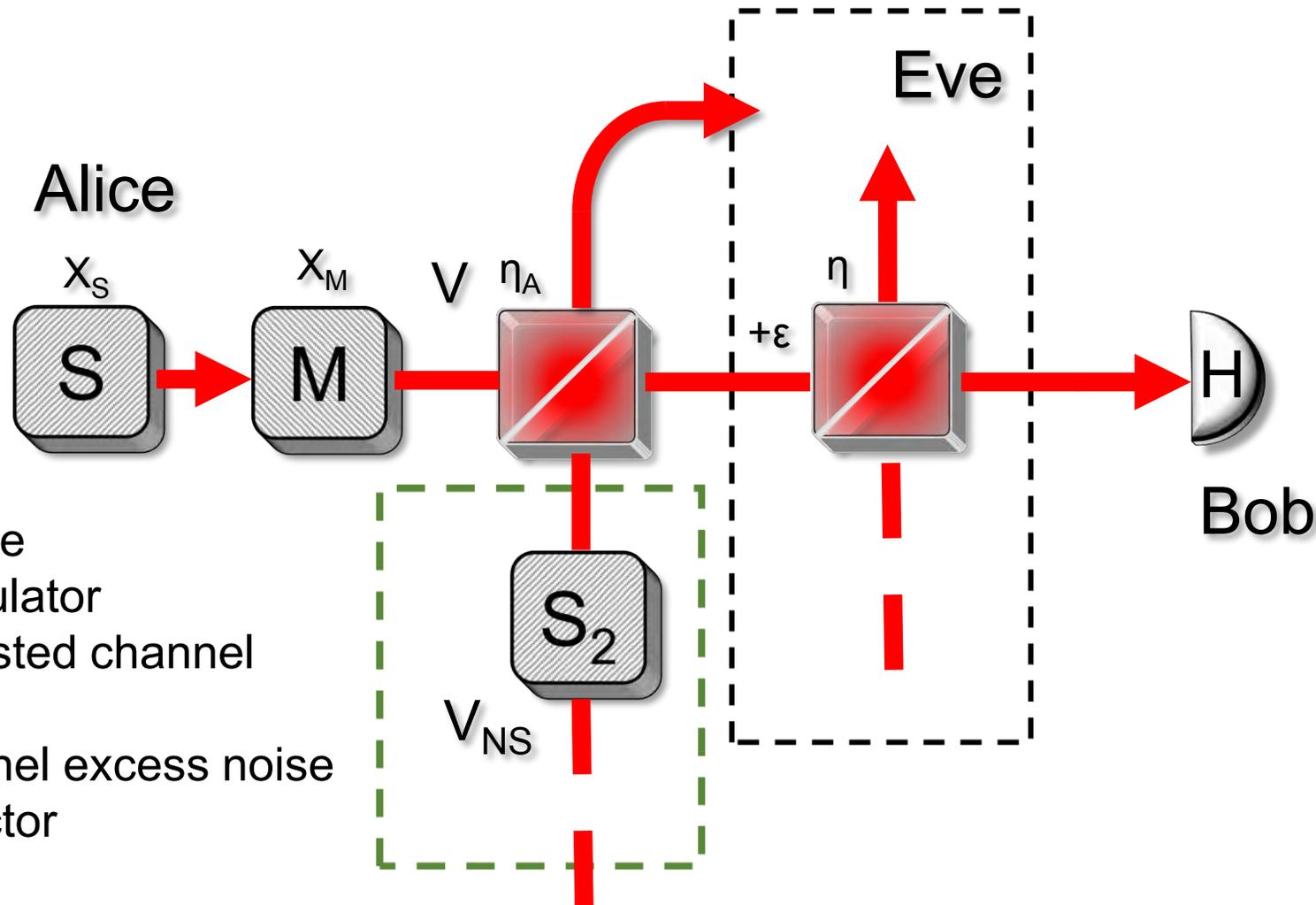


Side channels undermine the tolerance of the protocol to the channel noise ϵ .

Noise infusion on the receiver side leads to security break !

Decoupling of the type-A side channel

Leakage on the sender side

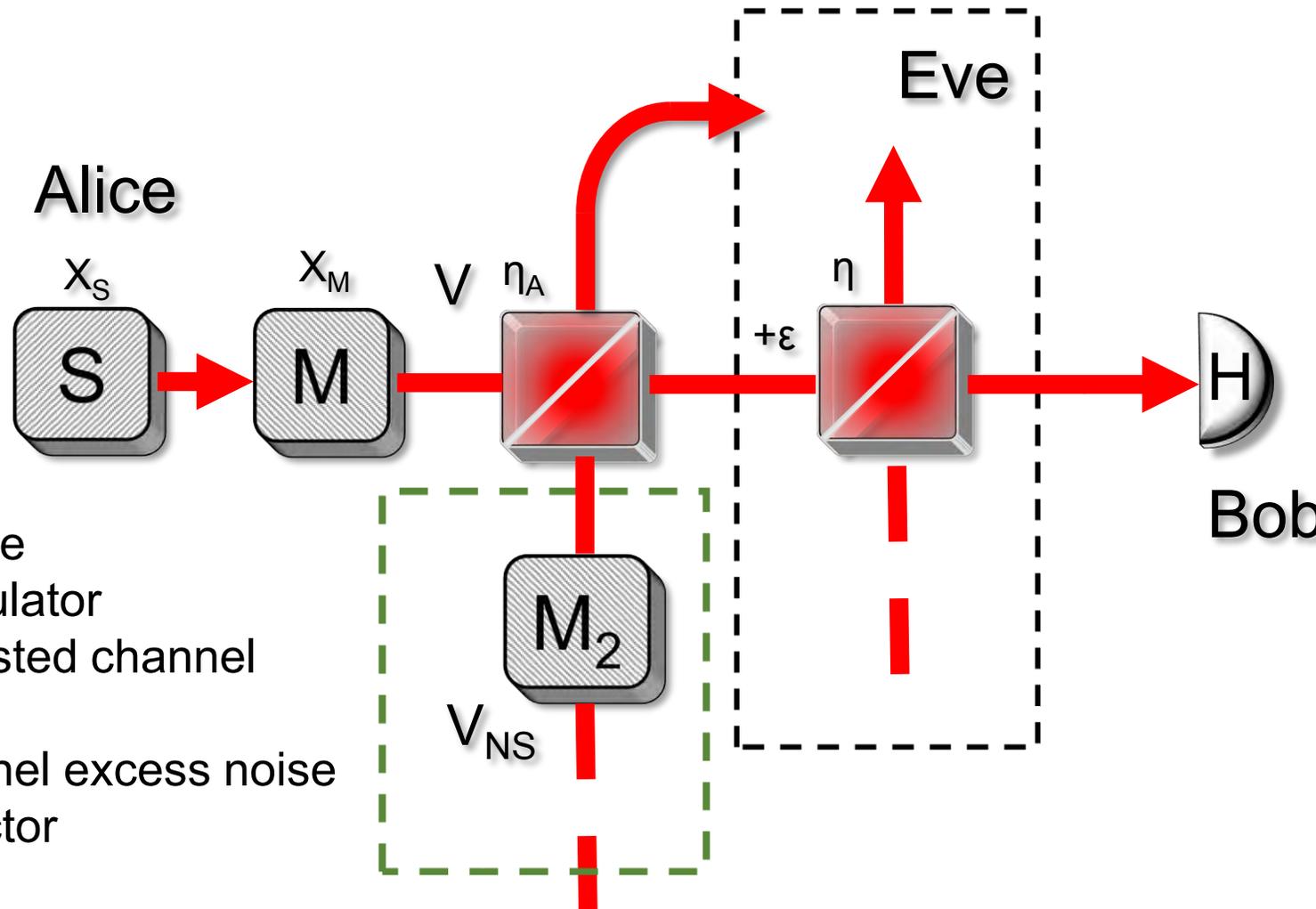


S - source
M - modulator
 η - untrusted channel loss
 ε - channel excess noise
H - detector

Alice replaces vacuum input of the side channel with the source of a Gaussian thermal noise (noise is unknown by definition). Holevo bound is reduced, but mutual information between trusted parties is reduced as well.

Decoupling of the type-A side channel

Leakage on the sender side



S - source
M - modulator
 η - untrusted channel loss
 ϵ - channel excess noise
H - detector

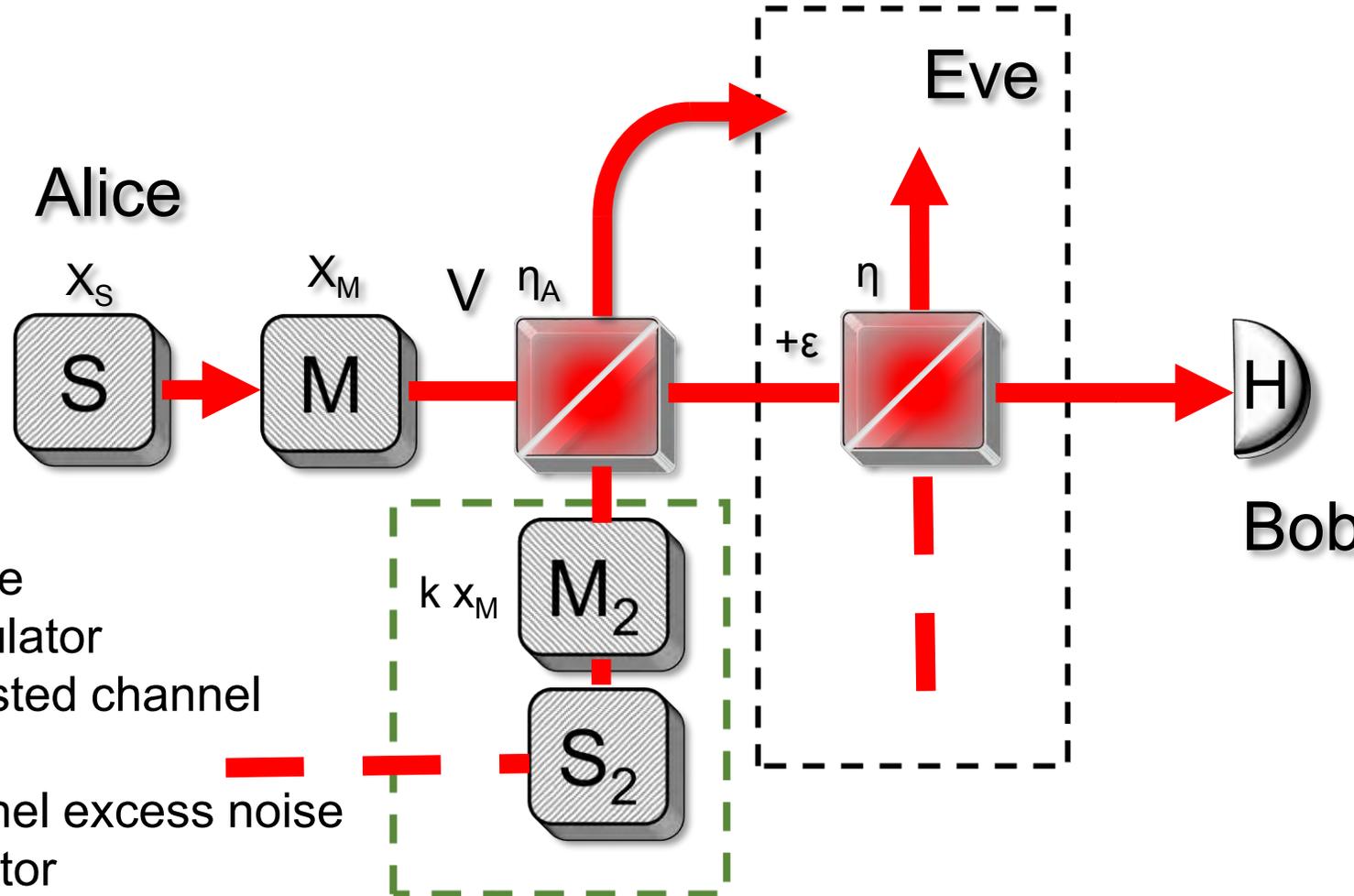
Alice replaces vacuum input of the side channel with modulator.

The modulation is independent from the main modulation but contributes to the trusted-side data and to correlation with the receiver.

Mutual information between trusted parties is increased, but the information leakage from the main channel is increased as well.

Decoupling of the type-A side channel

Leakage on the sender side



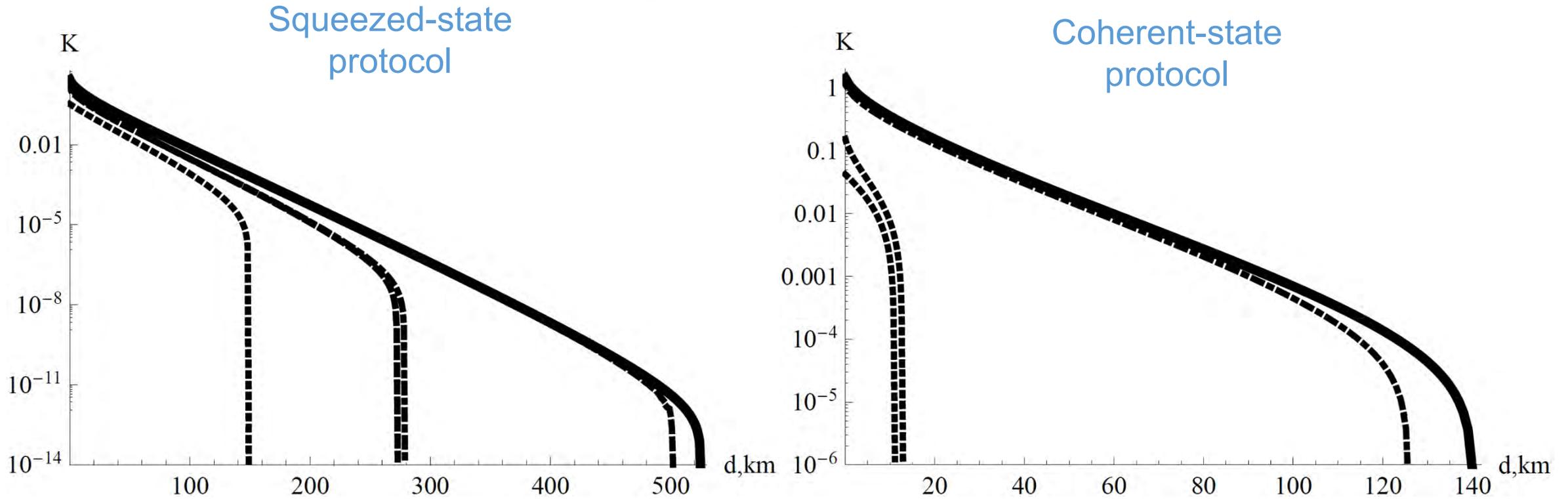
Alice replaces vacuum input of side channel with modulator. The modulation is correlated to the main modulation performed on the signal, and optionally squeezed.

$$k_{opt} = \sqrt{(1 - \eta_A) / \eta_A}$$

Optimal correlated modulation shifts side-channel attack from the modulated signal to signal state before modulation. Optimizing the parameters allows to **completely decouple and decorrelate** side channel output from the signal mode.

Lower bound on the key rate

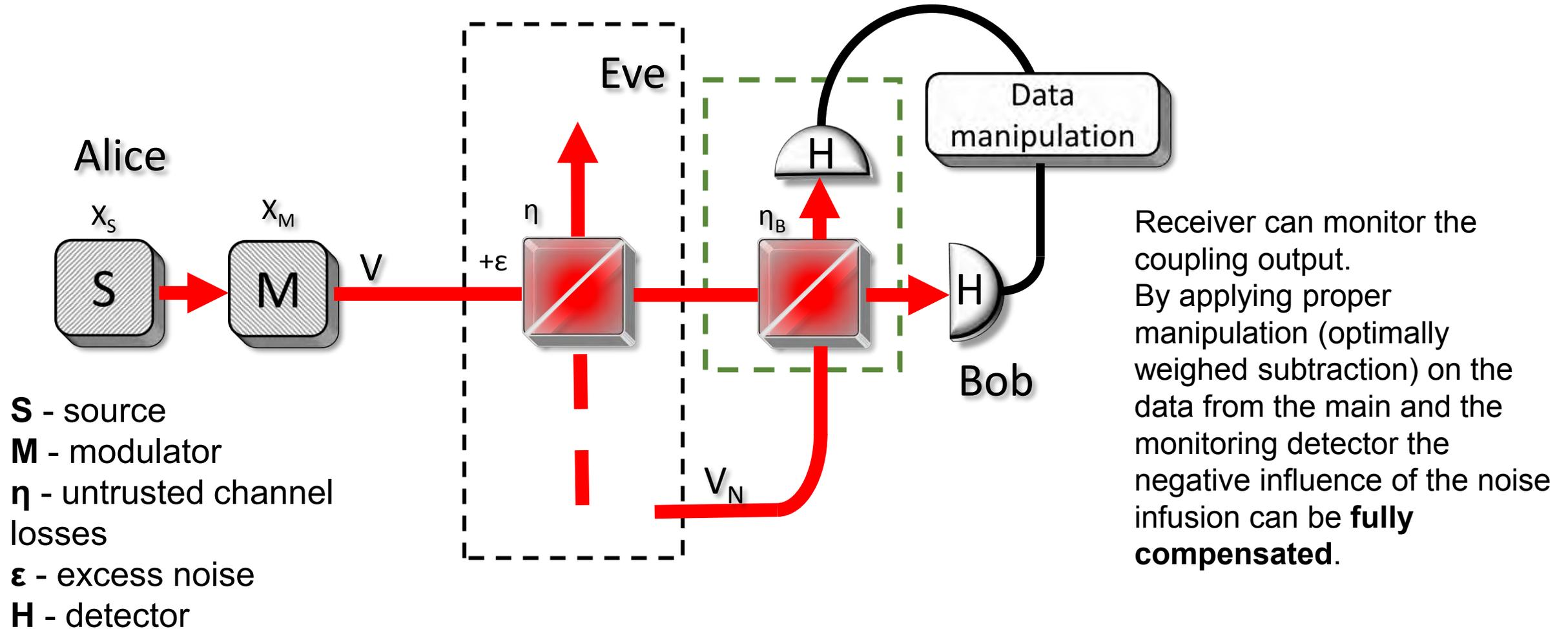
Leakage on the sender side



Channel losses: -0.2dB/km , Reconciliation algorithm efficiency: 95%, Untrusted channel excess noise: 5%, sender side side channel coupling: 0.4, Modulation variance is optimized for given parameters.

Decoupling of the type-B side channel

Noise infusion on the receiver side

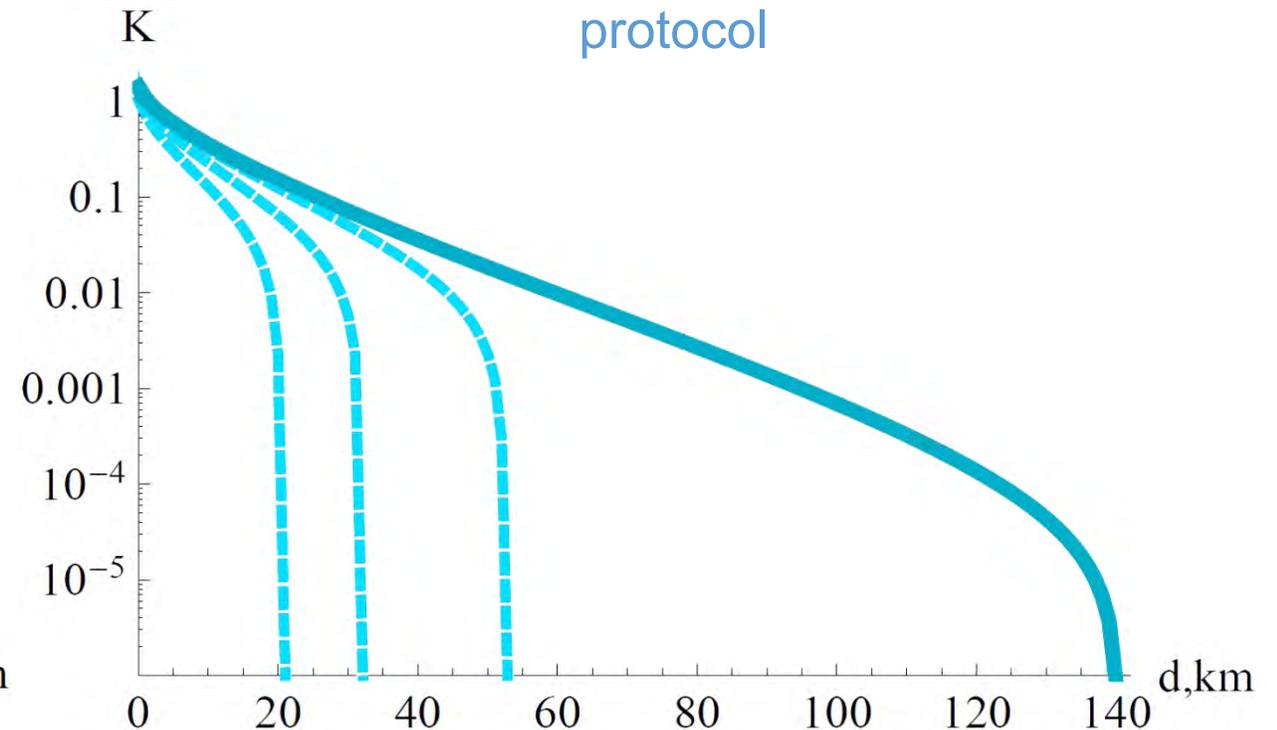
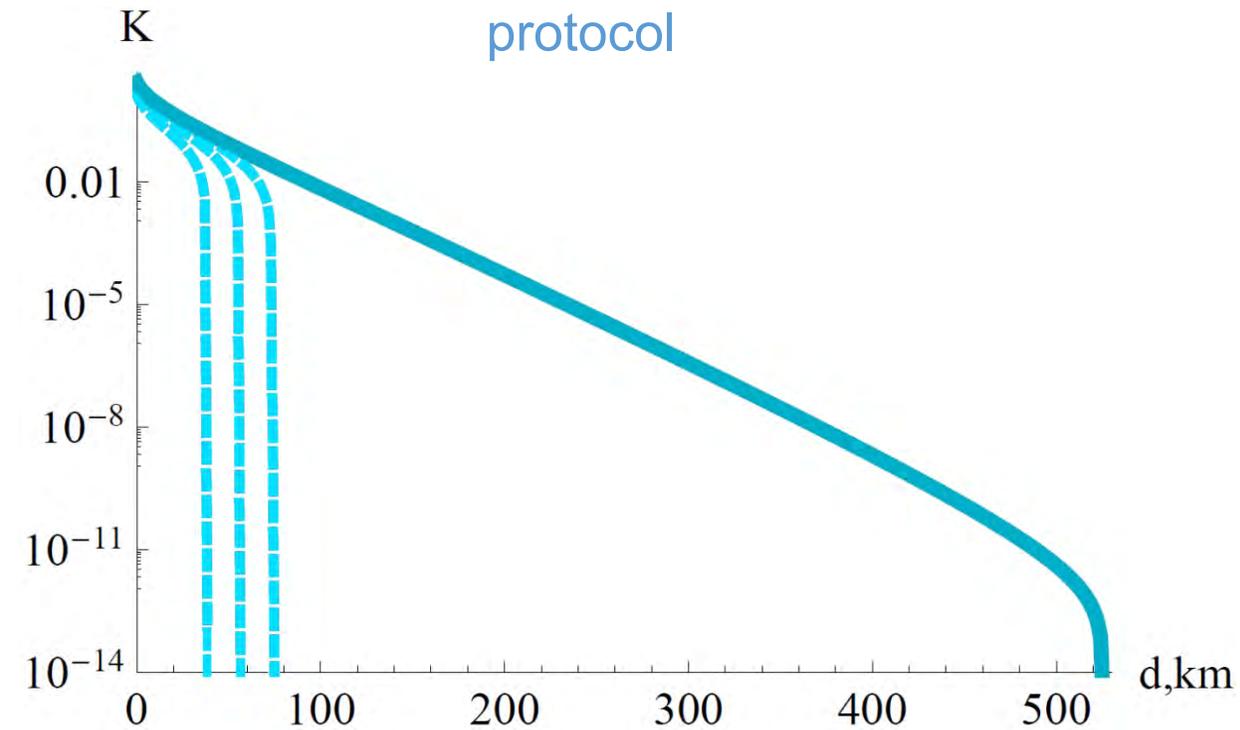


Lower bound on the key rate

Noise infusion on the receiver side

Squeezed-state
protocol

Coherent-state
protocol



Channel losses: -0.2dB/km , Reconciliation algorithm efficiency: 95%, Untrusted channel excess noise: 5%, Receiver side noise infusing side channel coupling: 0.5, 0.7, 0.9, Infused noise variance: 1.05

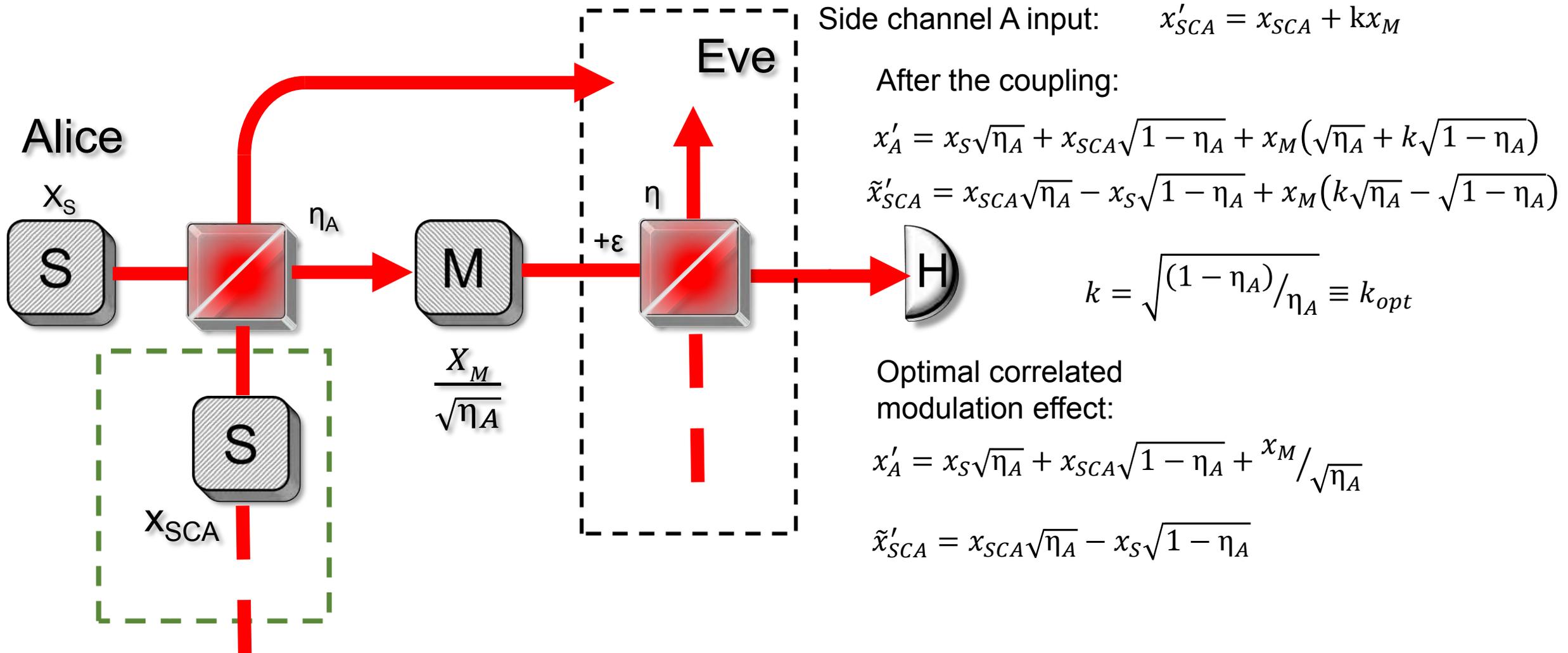
Summary

- Information leakage on the sender side decreases the key rate and increases sensitivity of the protocol to the channel noise. However by applying additional modulation and optionally squeezing to the input of side channel the negative impact can be decreased or even eliminated completely.
- Noise infusion on sender side can completely break the security of the protocols even upon pure channel loss. Nevertheless introduction of proper monitoring on the output of side channel allows to fully compensate the negative influence.

Thank you for your attention!

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Sender-side side channel decoupling



Receiver-side side channel decoupling

After noise infusion:

$$x'_B = x_B\sqrt{\eta_B} + x_{SCB}\sqrt{1 - \eta_B}$$
$$x'_{SCB} = -x_B\sqrt{1 - \eta_B} + x_{SCB}\sqrt{\eta_B}$$

Weighed subtraction:

$$\Delta x = g x'_B - g' x'_{SCB}$$
$$g = \sqrt{\eta_B} \quad g' = \sqrt{1 - \eta_B}$$

$$\Delta x = x_B$$