Introduction

We present a graph-based method of reasoning with if-then rules describing dependencies between graded attributes. The rules are usually written as expressions $A \Rightarrow B$, where A and B are fuzzy sets of attributes, and have two basic interpretations:

- :: Attribute implications in formal fuzzy context: "If it is (very) true that an object has all attributes from A, then it has also all attributes from B."
- :: Fuzzy functional dependencies in similarity-based relational databases: "For any two objects (rows): if they have (very) similar values on attributes from A then they have similar values on attributes from B''.

What those two interpretations have in common?

- :: same syntax (although different interpretation)
- :: same notion of semantic entailment
- :: single inference system can be used

Armstrong-like axiomatic system (L finite, Y set of attributes, $A, B, C, D \in \mathbf{L}^Y$):

(Ax)	$infer \ A \cup B \Rightarrow B$
(Cut)	from $A \Rightarrow B$ and $B \cup C \Rightarrow D$ infer $A \cup C \Rightarrow D$
(Mul)	from $A \Rightarrow B$ infer $c^* \otimes A \Rightarrow c^* \otimes B$,

Two kind of completeness can be shown:

- :: Ordinary completeness
 - Theory T set of if-then rules
 - $A \Rightarrow B$ is provable from T iff $A \Rightarrow B$ semantically follows from T (in degree 1)
- :: Graded (Pavelka-style) completeness
 - Theory T fuzzy set of if-then rules
 - A degree to which $A \Rightarrow B$ is provable from T equals to degree to which $A \Rightarrow B$ semantically follows from T.

Maier (2) proposed graph-based approach in order to normalize proof for reasoning with (ordinary) functional dependencies in relational databases. Can this approach be extended for graded if-then rules? YES.



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Lucie Urbanova, Vilem Vychodil: Derivation Digraphs for Graded If-Then Rules

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Derivation Digraphs for graded if-then rules

:: acyclic digraphs

- :: vertices are attributes from Y
- :: arcs correspond to if-then rules used from an input theory
- :: construction:

every set of unconnected vertices is derivation diagram new vertex for attribute y is added to graph if the validity of y resulting from T is strictly higher then the degree yhas in the current state of graph, formally:

T-based L*-derivation diagram (DAG): Let $\mathbf{L}^* = \langle L, \wedge, \vee, \otimes, \rightarrow, *, 0, 1 \rangle$ be be complete residuated lattice with a truth-stressing hedge, T be a set of graded if-then rules over Y.

- 1) Any $\mathbf{D} = \langle V, \emptyset \rangle$ such that $\emptyset \neq V \subseteq Y \times L$ is a *T*-based L^{*}-derivation DAG;
- 2) If $\mathbf{D} = \langle V, A \rangle$ is a *T*-based L*-derivation DAG and there are $E \Rightarrow F \in T$, attribute $y \in Y$, and vertices $\langle y_1, a_1 \rangle \in V, \ldots, \langle y_k, a_k \rangle \in V$ such that for
 - $s_0 = \bigwedge \{ E(y) \to 0 \mid y \in Y \text{ and } y \notin \{y_1, \dots, y_k\} \},\$ $s_1 = \bigwedge \{ E(y_i) \to a_i \mid i = 1, \dots, k \},$
 - $m = \bigvee \{ a \in L \, | \, \langle y, a \rangle \in V \},$ $d = ((s_0 \wedge s_1)^* \otimes F(y)) \vee m,$

we have d > m > 0, then $\mathbf{D'} = \langle V', A' \rangle$, where

 $V' = V \cup \{\langle y, d \rangle\},\$ $A' = A \cup \{ \langle \langle y_i, a_i \rangle, \langle y, d \rangle \rangle \mid i = 1, \dots, k \},$

is a T-based L*-derivation DAG.

One step in construction of DAG

- :: L*: finite linearly ordered Łukasiewicz algebra with $L = \{0, 0.1, \dots, 0.9, 1\} \subseteq [0, 1], \land \text{ and } \lor \text{ usual minima and } \}$ maxima, \otimes , \rightarrow Łukasiewicz operations, * identity.
- :: set of initial vertices: $\{\langle y_1, 0.7 \rangle, \langle y_2, 0.7 \rangle\}$ theory: $T = \{\dots, \{0.7/y_1, 0.8/y_2\} \Rightarrow \{0.6/y_3\}, \dots\}$

:: adding vertex $\langle 0.5, y_3 \rangle$



 $s_0 = 1$ $s_1 = 0.9$ m = 0 $d = (0.9 \otimes 0.6)^*$ = 0.5

Completeness

For T-based L*-derivation DAG D we have: :: yield of **D** on attribute y: $\mathbf{D}(y) = \bigvee \{a \in L \mid \langle y, a \rangle \in V\}$

:: D is called a T-based L*-derivation DAG for $E \Rightarrow F$ if $\{\langle y, E(y) \rangle \mid y \in Y \text{ and } E(y) > 0\}$ is the set of initial vertices of **D** and $\mathbf{D}(y) \ge F(y)$ for all $y \in Y$.

Results

Let T be a theory (set of if-then rules) :: If $T \vdash A \Rightarrow B$ (using (Ax),(Cut),(Mul)), then there is a *T*-based L*-derivation DAG for $A \Rightarrow B$.

:: If there is a T-based L*-derivation DAG for $A \Rightarrow B$, then $T \vdash A \Rightarrow B.$

:: If L is finite, then $A \Rightarrow B$ semantically follows from theory T in degree 1 iff there is a T-based L*-derivation DAG for $A \Rightarrow B.$

:: If L is finite, then $||A \Rightarrow B||_T$ is the greatest degree $a \in L$ such that there is a T-based L*-derivation DAG for $A \Rightarrow a \otimes B$.

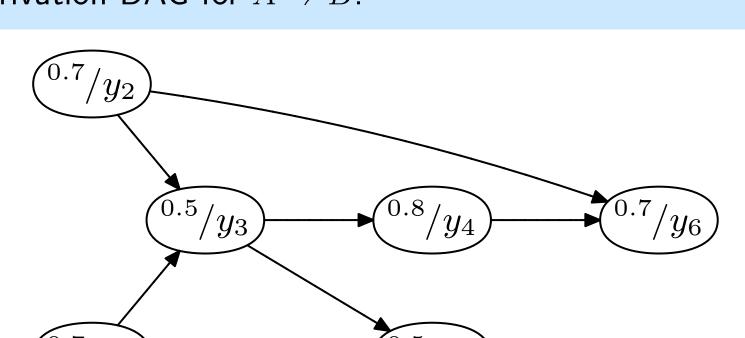
Example

We demonstrate here how to obtain a proof from a T-based L^* -derivation DAG. Assume we have a theory T:

$$T = \{\{{}^{0.7}/y_3\} \Rightarrow \{{}^{1}/y_4\}, \{{}^{0.7}/y_1, {}^{0.8}/y_2\} \Rightarrow \{{}^{0.6}/y_3\} \Rightarrow \{{}^{0.6}/y_1, {}^{0.9}/y_3\} \Rightarrow \{{}^{0.9}/y_5, {}^{0.1}/y_6\}, \\ \{{}^{0.5}/y_2, {}^{0.6}/y_4\} \Rightarrow \{{}^{0.7}/y_6\}\}$$

and consider $A = \{{}^{0.7}/y_1, {}^{0.7}/y_2\}, B = \{{}^{0.6}/y_4, {}^{0.6}/y_6\}.$

Does $||A \Rightarrow B||_T = 1$? YES, since there is T-based **L***-derivation DAG for $A \Rightarrow B$:



:: set of initial vertices is equal to A

:: $B(y) \leq \mathbf{D}(y)$ for all $y \in B$

The following sequence is a part of the corresponding proof for $A \Rightarrow B$ using derivation rules (Ax), (Cut), (Mul) and derived rules projectivity and additivity.

Corresponding Proof

- 1) $\{ {}^{0.7}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.7}/y_1, {}^{0.7}/y_2 \}$ $2) \{ {}^{0.7}/y_1, {}^{0.8}/y_2 \} \Rightarrow \{ {}^{0.6}/y_3 \}$ $3) \{ {}^{0.6}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.5}/y_3 \}$ 4) $\{ {}^{0.7}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.6}/y_1, {}^{0.7}/y_2 \}$ $5) \{ {}^{0.7}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.5}/y_3 \}$ $6) \{ {}^{0.7}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.7}/y_1, {}^{0.7}/y_2, {}^{0.5}/y_3 \}$ $7) \{ {}^{0.7}/y_3 \} \Rightarrow \{ {}^{1}/y_4 \}$ $8) \{ {}^{0.5}/y_3 \} \Rightarrow \{ {}^{0.8}/y_4 \}$

. $20) \{ {}^{0.7}/y_1, {}^{0.7}/y_2 \} \Rightarrow \{ {}^{0.6}/y_4, {}^{0.6}/y_6 \}$

Conclusions

:: graph-based inference for graded if-then rules

:: degrees of semantic entailment of fuzzy attribute implication (fuzzy functional dependency) from a theory fuzzy set of fuzzy attribute implication (fuzzy functional dependencies) can be characterized by the existence of directed acyclic graphs

Future work

- :: further properties of diagram, complexity issues
- :: characterization of syntactic closures via DAG
- :: algorithm for computing closures

References





(Ax) $\in T$ (Mul) on 2) (Ax)(Cut) (Add) $\in T$ (Mul) on 7)

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